Online Anomaly Detection With Bandwidth Optimized Hierarchical Kernel Density Estimators

Mine Kerpicci, Huseyin Ozkan, and Suleyman Serdar Kozat

Abstract—We propose a novel unsupervised anomaly detection algorithm that can work for sequential data from any complex distribution in a truly online framework with mathematically proven strong performance guarantees. First, a partitioning tree is constructed to generate a doubly exponentially large hierarchical class of observation space partitions, and every partition region trains an online kernel density estimator (KDE) with its own unique dynamical bandwidth. At each time, the proposed algorithm optimally combines the class estimators to sequentially produce the final density estimation. We mathematically prove that the proposed algorithm learns the optimal partition with kernel bandwidths that are optimized in both region-specific and time-varying manner. The estimated density is then compared with a data-adaptive threshold to detect anomalies. Overall, the computational complexity is only linear in both the tree depth and data length. In our experiments, we observe significant improvements in anomaly detection accuracy compared with the state-of-the-art techniques.

Index Terms—Anomaly detection, bandwidth selection, kernel density estimation, online, regret analysis.

I. INTRODUCTION

Anomaly detection, as an unsupervised one-class machine learning problem, finds use in various applications ranging from cyber security [1], surveillance [2], failure detection [3], novelty detection [4], and data imputation [5]. We study this problem for sequential data in a truly online setting and propose a novel algorithm that observes \(x_t \in \mathbb{R}^d\) at each time \(t\) and decides whether \(x_t\) is anomalous based on the past observations \(\{x_1, x_2, \ldots, x_{t-1}\}\). In our framework, each instance \(x_t\) is processed only once without being stored and no label information is required. Therefore, the proposed algorithm is online and unsupervised, and appropriate for real-time anomaly detection applications. One of the main difficulties in anomaly detection problems is typically the extreme rarity of anomalies and their unpredictable nature [6], which differentiates the anomaly detection from the binary classification. In the binary classification problem, one has observations from both classes. However, in anomaly detection, there are no or extreme rare observations from the anomaly class. This essentially poses a one-class classification problem where it is reasonable to assume uniformity for anomalies [7]. Then, the Neyman–Pearson (NP) test [8] corresponds to comparing the density \(f(x_t)\) with an appropriate threshold [9], [10] that maximizes the detection power at a desired false alarm rate. This in turn amounts to density estimation and learning the desired threshold. Therefore, a common approach in the literature [6] consists of two steps, which are assigning a density score to the observed data instance and, then, comparing this score with a threshold [6]. Scores below the threshold indicate anomaly.

Our technique is also twofold. We first estimate the probability density of the data in a nonparametric manner without any model assumption and then find the anomalies by thresholding the estimated density for each instance.

Our density estimator is regret-optimal [11], where the loss is the accumulated negative log likelihood, against a large class of estimators. To be more precise, our density estimator asymptotically performs at least as well as the best class estimator in terms of the data likelihood. To obtain an estimator in this class, we use a kernel density estimator (KDE) [12] with a different bandwidth parameter in each region of a given partition of the observation space. Hence, the class of estimators is restricted to KDEs defined on the tree where varying the bandwidth parameter in each region and also varying the partition itself yield the aforementioned class. The proposed algorithm learns the optimal partitioning (that can be defined over the introduced tree) of the observation space and also learns the optimal region-specific bandwidth parameters for the optimal partition even in a time-varying manner. Namely, the optimal bandwidth of the KDE that we use for each region in the optimal partition might well be different as well as nonstationary (the distribution that the data instances are drawn/coming from might be changing in time, i.e., might be time-varying).

The introduced density estimation for anomaly detection is essentially a mixture of expert-based [13] approach. These experts are “piecewise” KDEs, “pieces” of which are generated by a hierarchical observation space partitioning via a depth-\(D\) binary tree. Pieces correspond to the partition regions, and for each region, the optimal time-varying KDE bandwidth from a finite set with cardinality \(k\) is separately learned. Our algorithm produces its density estimation by combining

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these expert estimations. The proposed combination is based on a time-varying weighting that is used to mitigate overfitting/overtraining issues and to asymptotically pick the optimal partition. The purpose in hierarchically organizing the KDEs is to achieve computational efficiency while still maintaining an exponentially large competition class of estimators. Specifically, the introduced tree-based hierarchical structure provides us a class of size approximately $(1.5)^{2D}$, whereas the overall processing complexity in our framework scales linearly only with $O(TDk)$, where $T$ is the number of processed instances.

Our initial findings belonging to the use of KDE with only a single partitioning along with the bandwidth selection have been presented in a national conference [14]. On the other hand, in this article, we complete our study comprehensively and provide our final greatly detailed experimental results. In particular, in this article, we generalize to a hierarchically structured competing class of doubly exponentially many partitionings (i.e., density estimators) with our proposed bandwidth selection method while additionally achieving strong mathematical performance guarantees, which therefore yields significantly higher anomaly detection performance as presented in our extensive experiments.

The result is a novel computationally efficient combination of hierarchically organized piecewise online KDEs with bandwidths that are optimally varying in time and space. We mathematically prove and guarantee the convergence of the density estimation performance of our algorithm to the best piecewise constructed class estimator in terms of the data likelihood. Finally, we decide whether the observed data instance is anomalous or not by comparing the estimated probability with a threshold. In our approach, an optimal NP test can be achieved with an appropriate threshold choice for a constant false alarm rate. Alternatively, we provide an algorithm to adaptively learn the optimal threshold in terms of the detection accuracy when the true labels are available. Also, we demonstrate significant performance gains by our algorithm in comparison to state-of-the-art methods through extensive set of experiments.

A. Related Work

Kernel-based techniques are extensively used in the machine learning literature and in particular for density estimation [15] due to their nonparametric high modeling power. Nevertheless, kernel bandwidth selection is one of the most critical issues in these approaches [16] since the bandwidth choice significantly affects the performance of the density estimator. There are two main approaches to handle this issue, which are based on cross validation and plug-in [17]. Cross-validation-based approaches mostly depend on the leave-one-out scheme to find the bandwidth [17]. Plug-in-based methods are based on the minimization of asymptotic mean integrated squared error (AMISE) [17], where pilot bandwidths are introduced for an iterative solution. Another method is proposed in [18], which pointwise varies the kernel bandwidth based on the minimization of error such as AMISE. Although these techniques and others such as [19] and [20] are impressive in their own batch processing setting, they suffer from one or both of the following two issues: 1) they use one global time-invariant bandwidth for the complete observation space and time span and 2) their computational complexity is prohibitively large, preventing their use in our sequential setting. There are also online KDE methods, which use compression techniques to adapt the kernel method to an online framework, with time-varying bandwidth optimization as in [21]–[23]. These methods aim to reduce the number of stored data instances to lower complexity. However, this compression approach requires iterative steps to update the compressed model and bandwidth at each time, which is prohibitively costly in an online framework. An online adapted version of AMISE-optimal bandwidth approach is proposed in [21], which requires an iterative plug-in process for each observed data instance. The compressed model of samples is used in [22] to find the AMISE-optimal bandwidth at each time after the model is updated with the observed data instance. However, they are not able to learn the local variations of data due to their bandwidth optimization approaches. Our method, however, rigorously addresses all of these issues by the introduced highly adaptive hierarchical KDEs and their locally optimized bandwidths with low computational complexity.

There exist numerous methods [9], [24]–[26] that solve the anomaly detection problem in the nonsequential batch setting. For example, Zhao and Saligrama [9] and Dasarathy [25] computed pairwise cross $k$th distances between the observed instance and training samples for detecting an anomaly. Another method is proposed in [26], which is less sensitive to the neighborhood parameter $k$ compared with the conventional distance-based techniques. Since these methods are all batch processors, repeated use of all of the available data is required that leads to too large memory and computational complexity. Among the NP-based anomaly detection techniques [7], [10], [27], minimum volume set approach is used in [10] for detecting anomalies. Also, local nearest neighbor distances are used for the anomalies in video observations in [7]. These are again both batch techniques for which an online extension is proposed in [27] that can only work when the data are Markovian. In contrast, our technique is truly online and model free, which is appropriate for real-time processing of challenging data sources.

Anomaly detection problem is solved in the online setting in [28]–[32]. For example, Xie et al. [29] proposed a $k$-nearest neighbor technique, with required training steps. Online training is introduced in [32] to adapt the batch algorithms, such as support vector machine (SVM) to real-time applications. Since these methods [29], [32] still require quadratic optimization, the computational running time significantly increases with the data length. Anomaly detection is considered as a classification problem in [30] and it is solved via a cost-sensitive approach, which requires true label information, whereas our technique is completely unsupervised. Raginsky et al. [28] and Filippone and Sanguinetti [31] estimated the probability density of the observed data and compared it with a threshold. For this, Raginsky et al. [28] assumed that the data distribution is from an exponential family. An information-theoretic approach is proposed in [31] to increase the detection performance when
the number of training data samples is small. These methods [28], [31] assume that the data distribution is Gaussian or a mixture of Gaussians and solves the anomaly detection problem with a parametric approach. Unlike these methods, we do not assume any parametric distribution shape for the observed sequential data, and hence, our technique can model any unknown arbitrarily complex distribution in a data-driven online manner with random Fourier features [33]. Kernel approximation with random Fourier features [33] is used for binary classification problem in [34] where the proposed algorithm learns (in online manner) the linear discriminant in the approximated and linear kernel space with online gradient descent (OGD) based on the true label information. Our method also uses random features for obtaining an online implementation but, on the other hand, solves the online anomaly detection problem by learning both the optimal partition in a hierarchical tree and the optimal bandwidths for all regions in this partition without requiring any true label information.

B. Contributions

Our contributions are as follows.

1) For the first time in the literature, we introduce a piecewise nonparametric KDE with a highly adaptive hierarchical tree structure. This addresses the bandwidth selection problem in a locally adaptive manner. Hence, we achieve a great modeling power for data of any arbitrarily complex distribution while successfully mitigating the overfitting by a carefully designed weighting scheme over the introduced class of experts.

2) We efficiently implement our kernel-based anomaly detection algorithm in an online setting. To this end, we use a hierarchical approach to represent and manage a very large class of estimators (1.5)D to obtain the final estimation with complexity only \( O(TDK) \) (\( T \): data size, \( D \): tree depth, and \( k \): cardinality of the bandwidth set). This is appropriate for real-time processing while achieving a highly powerful estimation power.

3) Our algorithm achieves the performance of the best estimator in class of estimators with a regret bound \( O(2^D/h) \) (\( D \): tree depth and \( h \): learning rate).

4) We illustrate significant performance improvements achieved by our algorithm with respect to the state of the art as well as the recently proposed methods through an extensive set of experiments over synthetic and real data.

C. Organization

We first define our problem setting in Section II. In Section III, we introduce the hierarchical partitioning of the observation space. Section IV presents our combination of hierarchically organized online KDEs for our density estimation. Section V introduces our optimal time-varying thresholding for the final anomaly detection. In Section VI, we demonstrate our performance improvements on artificial and real data sets. We conclude with final remarks in Section VII.

II. Problem Description

We sequentially observe \( x_t \in \mathbb{R}^d \) at each time \( t \). Our aim is to decide whether the observed data \( x_t \) is anomalous or not based on the past observations, i.e., \( x_1, x_2, \ldots, x_{t-1} \). For this purpose, we introduce a two-step approach for anomaly detection where we first sequentially estimate the probability density of \( x_t \) by using previously observed data sequence \( x_1^{t-1} = \{x_i\}_{i=1}^{t-1} \) as \( f_t(x_t) \). Then, we compare the estimated probability \( f_t(x_t) \) with a threshold to decide whether \( x_t \) is anomalous or not. We declare a decision \( d_t \) at each time \( t \) as

\[
\hat{d}_t = \begin{cases} 
+1, & f_t(x_t) < \varsigma \text{ (anomalous)} \\
-1, & f_t(x_t) \geq \varsigma \text{ (normal)} 
\end{cases}
\]

where \( \varsigma \) is the threshold. We also provide an algorithm that updates the threshold \( \varsigma_t \) in time to reach the best choice in terms of the detection accuracy when the true label \( d_t \) is available. We emphasize that our framework is completely unsupervised, i.e., no knowledge of the true label is required, but if the true label is available for an instance, then it is incorporated.

We assume that the observation sequence \( \{x_t\}_{t=1}^T \) is generated from an independent and identically distributed (i.i.d.) source where \( S \) is a bounded compact domain. In order to estimate the probability density of \( x_t \), we introduce a novel algorithm based on the context tree weighting method [35], [36] in which we construct a class of density estimators and asymptotically achieve—at least—the performance of the best class estimator.

Our goal is first to sequentially estimate the density of each instance \( x_t \) causally in time by combining the class estimators that we construct with hierarchical partitioning of a tree and asymptotically achieve at least the performance of the best (in terms of the accumulated likelihoods) class estimator. This combination can be seen as a mixture of experts with each expert being a class estimator. Second, we optimally (in terms of the anomaly detection accuracy) threshold our estimated density for anomaly detection.

We use log-loss, i.e., negative log likelihood, as our pointwise loss function: \( l(f_t(x_t)) = -\log(f_t(x_t)) \) since this loss successfully encodes the data likelihood. Then, our goal of asymptotically performing at least as well as the best density estimator in the introduced class is equivalent to the minimization of the following regret \( R(T) \) [11], resulting in a regret-optimal final estimator:

\[
R(T) = \sum_{t=1}^{T} (-\log(f_t(x_t)))
\]

\[
-\min_{P_t} \left\{ \sum_{t=1}^{T} (-\log(f_t^P(x_t))) \right\}
\]

1In this article, all vectors are column vectors, and they are denoted by boldface lower case letters. For a vector \( w \), \( w^\top \) is its transpose. Also, for a scalar \( x \), \( |x| \) represents its absolute value, and for a set \( X \), \( |X| \) represents the cardinality of the set. The time index is given as subscript, and a sequence of \( \{x_t\}_{t=1}^{N} \) is represented as \( x^N \). Also, \( 1_{\{\cdot\}} \) is the indicator function returning 1 if its argument condition holds and returning 0 otherwise.
where $T$ is the length of the sequence $x^T = \{x_t\}_{t=1}^T$ and the regret is the performance difference between our proposed estimator $f_t$ and the best class estimator. Here, $\hat{f}^P_i$ is the piecewise constructed KDE for the partition $P_i$. In Section IV, we achieve the introduced goal by proving that this regret is bounded by a term that is convergent to 0 after normalization with $T$.

We compare the estimated density $f_t(x_t)$ with a time-varying threshold $\zeta$ to optimally decide (in terms of the detection accuracy) whether the observed data $x_t$ is anomalous as in (1). For this purpose, by using the logistic function $r_t = \log(1 + \exp(-\zeta - f_t(x_t)))d_t$, we update the threshold to maximize the detection power when the true label $d_t \in \{-1, +1\}$ is available. If the true label is never available, then one can straightforwardly adjust the threshold to bound the false alarm rate in a similar manner. Recall that the proposed algorithm is unsupervised and can certainly operate without requiring the label information.

### III. TREE CONSTRUCTION

In this section, we present the hierarchical partitioning tree to divide the observation space into disjoint regions and construct a class of density estimators.

In the literature, trees are used in various ways to improve the performance of the learning methods. The tree structure is used in a change detection problem in [37]. The use of randomized binary trees in the ensembles for outlier detection is pointed out in [38]. Moreover, a binary splitting scheme for the change detection problem is proposed in [39]. In our framework, we use a tree structure to construct a class of density estimators localized to various regions and combine their decisions to achieve better performance.

A hierarchical partitioning tree of depth-$D$ is constructed to partition the observation space $S$ into various regions such that each tree node $\theta^2$ corresponds to a specific region in $S$. We start splitting from the middle of the first dimension of the observation space. At each depth, we divide the observation space from the middle of the latter dimension. For an illustration, in Fig. 1(a), we consider two dimensional bounded $x_t \in \mathbb{R}^2$, i.e., $x_t = \{x_{t,1}, x_{t,2}\}$ and $|x_{t,1}| < A$, and use a depth-2 tree. There exist $2^{D+1} - 1$ nodes for a depth-$D$ tree in general, and each node region is the union of the regions of its left and right children. For example, the children of the root node $\theta_0$ are defined as $\theta_{ol} = [-A, 0] \times [-A, A]$ and $\theta_{or} = [0, A] \times [-A, A]$ in Fig. 1(a). The root node $\theta_0$ corresponds to the complete observation space such that $\theta_0 = \theta_{ol} \cup \theta_{or} = S = [-A, A] \times [-A, A]$.

**Remark:** Note that there is no limit for the choice of depth $D$ since partitioning is used to construct the disjoint regions, and increasing the number $2^{D+1} - 1$ of nodes increases (with sufficient data) the performance of KDE. However, the depth $D$ should be chosen carefully in practice with limited data. In this study, we have conducted separate preliminary experiments in each data set we used and observed that small $D$ around $D = 3$ or $D = 4$ generally provides a highly good performance. Hence, we opted not to further optimize $D$ and set it to $D = 3$ across our data sets. On the other hand, we emphasize that the required number of partition regions scales exponentially with the feature dimension $d$ if one desires a specific-level granularity in partitioning, and it scales linearly if one desires to split each dimension at least once. This poor scaling with the ambient dimension for the desired degree of fine partitioning appears to be a limitation of our presented technique. To address this limitation and handle high dimension, we apply principal component analysis (PCA) to the data and reduce its dimensionality to $d = 3$ prior to depth selection with a preliminary experiment. Hence, with PCA, our technique can still be applied successfully in high dimension especially when the covariance structure in the data allows good compression, which we strongly verify in our experiments. Note that PCA is only used for partitioning and constructing the tree, after which the instances traverse over tree with their full dimension. Namely, for an instance, we decide the tree nodes based on the reduced dimension, but the KDE at each node runs on the full dimension.

We combine $n_i$ disjoint regions on the tree to form a partition $P_i$ such that $P_i = \{\theta_{i,1}, \theta_{i,2}, \ldots, \theta_{i,n_i}\}$ with $\theta_{i,k} \cap \theta_{i,j} = \emptyset$.
if \( k \neq j \) and \( \bigcup_{j=1}^{n_p} \theta_{i,j} = S \). Examples of such partitions in the case of a depth-2 tree are shown in Fig. 1(b). In general, we construct \( n_p \approx (1.5)^{2D} \) [35] partitions for depth-\( D \) tree. Note that each of these partitions can be obtained by collecting the regions of the leaves after an appropriate pruning. One can obtain five different partitions for \( D = 2 \) as in Fig. 1(b). We assign a KDE to each region \( \theta_{i,j} \) in each partition \( \mathcal{P}_i \) such that the KDE is trained with only the instances falling in the corresponding region. Note that each partition essentially collects \( n_i \) many region-specific KDEs yielding a “piecewise” constructed KDE. Consequently, the set of \((1.5)^{2D}\) partitions from all possible prunings of the tree defines a class of piecewise constructed KDEs.

IV. PROPOSED NOVEL DENSITY ESTIMATOR

In this section, we propose a novel density estimator based on the hierarchical partitioning of the observation space using a binary tree. Our approach is to sequentially train nonparametric KDEs [12] at each node of the introduced tree by instances falling in the corresponding region while separately optimizing the bandwidth parameter [16] for each region. This process, and thus the proposed algorithm, is online and sequential with no storing.

A. Bandwidth Optimized Piecewise Online KDE

The main building block of our anomaly detection algorithm is the online bandwidth optimized KDE employed in regions of the observation space, i.e., at each node of the defined tree. In this section, our discussion about this building block is confined to a single region from the observation space. In Section IV-B, we introduce the combination of KDEs from all possible regions.

1) Conventional KDE: We use the KDE method at each node \( \theta_k \) to estimate the local density in the corresponding region of the observation space based on the instances falling in that region. Let \( x \) be any point in \( \theta_k \), i.e., \( x \in \theta_k \), at time \( t \) before observing \( x_t \) along with the past data \( \{x_i\}_{i=1}^{t-1} = x^{t-1} \). Then, KDE is obtained as

\[
\tilde{f}_t^k(x; \delta) = \frac{1}{\pi^{d/2} \delta^d} \sum_{x \in \theta_k, x_t \in \mathcal{P}_i, 1 \leq r \leq l-1} K\left(\frac{x - x_r}{\delta}\right)
\]

(3)

where \( \tau_{t-1}^k = \sum_{i=1}^{l-1} \mathbb{1}_{\{x_t \in \mathcal{P}_i\}} \) is the number of instances falling in \( \theta_k \) until time \( t - 1 \), \( \delta \) is the bandwidth, and \( K(\cdot) \) is a kernel function [12]. In this work, we use the Gaussian kernel (or radial basis function) [33] \( K(x) = (1/(2\pi)^{d/2}) e^{-\|x\|^2/2} \). After inserting this into (3), we obtain

\[
\tilde{f}_t^k(x; g) = \frac{1}{\pi^{d/2} \delta^d} \sum_{x \in \theta_k, x_t \in \mathcal{P}_i, 1 \leq r \leq l-1} k(x, x_r; g)
\]

(4)

where \( k(x, x_r; g) = (1/N_k) \exp(-g\|x - x_r\|^2) \) with \( g = (1/2\delta^2) \) being the bandwidth parameter and \( N_k = (\pi/g)^{d/2} \) the normalization constant.

2) Online KDE: The summation in (4) requires all of the pairwise kernel evaluations between the test instance \( x \) and all the previous instances, which is computationally prohibitively complex. For an efficient kernel evaluation, we explicitly map the data points \((x, x_r)\) in the kernel evaluation into a high-dimensional kernel feature space with an explicit and randomized “compact” (cf. the remark below for compactness) mapping \( \phi: \mathbb{R}^d \to \mathbb{R}^m \) such that the kernel evaluation can be arbitrarily well approximated as an inner product of the mapped points [33] \( k(x, x_r; g) \approx \phi(x)\phi(x_r) \). In order to find such a mapping \( \phi(\cdot) \), we exploit the fact that Fourier transform of a symmetric shift-invariant and continuous positive-definite kernel represents a probability distribution [33], yielding

\[
k(x, x_r; g) = \int f_{w,b}(w, b; g) \times \left(\frac{g}{\pi}\right)^{d/2} \sqrt{2} \cos(w'x + b)
\]

\[
\times \left(\frac{g}{\pi}\right)^{d/2} \sqrt{2} \cos(w'x_r + b)dwdb
\]

(5)

where \( f_{w,b}(w, b; g) = (4\pi g)^{-d/2} e^{-4\pi g/(\|w\|^2/k)} \times (1/2\pi1[0 \leq b < 2\pi]) \) is a valid probability density function obtained as a result of the Fourier transform of the kernel being used together with an auxiliary uniform randomization over \( b \) [33]. Then, by approximating the expectation result for the kernel evaluation in (6) through a sample mean (due to the law of large numbers), we obtain the final approximation with the mapping function

\[
\phi(x; g) \triangleq \sqrt{\frac{2}{m}} \left(\frac{g}{\pi}\right)^{d/2} \times \left[ \cos(\sqrt{2gw'x + b_1}), \cos(\sqrt{2gw'x + b_2}), \ldots, \cos(\sqrt{2gw'x + b_m}) \right]'
\]

(7)

where \( (\{w_i, b_i\})_{i=1}^{m} \) is a random sample from \( f_{w,b}(w, b; (1/2)) \) and \( (w_i, b_i) \) are i.i.d. realizations [33]. Note that the algorithm that we introduce can be used with any symmetric shift-invariant and positive-definite kernel by only using the correct randomization \( f_{w,b}(w, b; (1/2)) \) (which is the Fourier transform of the kernel) and modifying \( \phi(x; g) \) accordingly [33].

Remark: Note that \( w \) and \( b \) form \( f_{w,b}(w, b; g) \), the probability density function in (6), such that the dimensions of \( w \) are independent scalar random variables and each has a Gaussian distribution with 0 mean and 2\( g \) variance, and \( b \) is a scalar random variable that is uniformly distributed between 0 and \( 2\pi \). In our method, we also aim to use adaptive bandwidth by changing \( g \) in time to improve the performance of the estimation. Therefore, we define our mapping function in (7) with \( \cos(\sqrt{2gw'(x + b)}) \) to make the sampling procedure of \( w \)'s independent of \( g \). This modification corresponds to drawing samples from \( f_{w,b}(w, b; (1/2)) \) such that the dimensions of \( w \)
have a Gaussian distribution with 0 mean and 1 variance, and $b$ is uniformly distributed between 0 and $2\pi$. Then, we multiply with $\sqrt{2\pi}$ in (7) to adjust for the desired covariance.

We obtain the density estimation in (4) with the sample mean approximation

$$
\hat{f}_t^\theta (x; g) = \frac{1}{t} \sum_{i=1}^{t} \phi(x; \theta_i) = \frac{1}{t} \sum_{i=1}^{t} \sum_{l \leq r < t} \phi(x; \theta_r) \phi_l \phi^* \tag{8}
$$

where

$$
\phi_l \phi^* \triangleq \frac{1}{t} \sum_{r \leq t} \phi(x; \theta_r)
$$

which can be implemented through sequential updates (after normalization with the total probability mass in $\theta_k$) as in line 11 in Algorithm 1.

Remark: This approximate form (but asymptotically exact form as $m \to \infty$) in (10) for the density estimation in (4) removes the pairwise kernel evaluations and allows us to evoke only a single dot product for the whole expression. This leads to an efficient online processing at the cost of the explicit and randomized mapping of each instance through $\phi(\cdot)$. We emphasize that this mapping is “compact” and the addition of mapping cost is fairly low; the quality of the approximation of the expectation through the sample mean in (10) improves at an exponential rate in $m$ due to Hoeffding’s inequality. Consequently, we obtain an online kernel density estimation at a negligible transformation cost.

3) Piecewise Online KDE: We point out that every pruning of our tree defines a partition of the observation space with regions attached to the leaves of the pruned tree, and there is an online KDE at each region. Following this observation, for a given partition $P = \{\theta_1, \theta_2, \ldots, \theta_N\}$, we combine the online KDEs $\hat{f}_i^\theta$’s at the partition regions $\theta_i$’s, i.e., pieces and form a piecewise online KDE $\hat{f}_t^P$ to model the density for the complete data. Recall that an online KDE is responsible for only those instances falling in its region, and hence, each online KDE must be normalized with the corresponding probability mass ($\tau_\theta / t - 1$) in its region $\theta_i$ while combining them. As a result, for each partition $P$ (there are approximately $(1.5)^{2D}$ many partitions defined over our tree), we define a piecewise online KDE $\hat{f}_t^P (x; g)$ as

$$
\hat{f}_t^P (x; g) = \sum \limits_{\theta \in P} \hat{f}_t^\theta (x; g) = \sum \limits_{\theta \in P} \phi(x; \theta) \phi_l \phi^* \tag{11}
$$

4) Bandwidth Optimized Piecewise Online KDE: In addition, our piecewise online KDE learns the optimal bandwidths separately for its constituent local estimators. For this purpose, we allow the bandwidths for the constituent local estimators to be different from region to region and also time-varying in line with the increasing data length. Hence, our technique is tuned to local variations of the data manifold as well as the rate of increase in the data size of the stream even in a nonstationary environment. In contrast, conventional approaches use one global KDE with a fixed time-invariant bandwidth that is certainly a strong limitation that we remove.

For this purpose, we use a finite set of bandwidths and run our online local KDEs in parallel with all bandwidth values in the set. The computational complexity of this parallel running in our approach is only $O(D)$ since the local KDEs of only those regions containing $x_i$ must be updated. In time, our approach identifies the optimal bandwidth $g^*$ in node $\theta_k$, which has the highest log likelihood (after scaling with $h$ that is exactly the log likelihood when $h = 1$)

$$
Z_t^{\theta_k} (g^*) = \max \limits_{g} \prod \limits_{j} \exp \{h \log \hat{f}_j^{\theta_k} (x_{ij}; g)\}.
$$

Note that the introduced optimality can be improved to any desired degree by using a larger finite bandwidth set $G = \{g_1, g_2, \ldots, g_{|G|}\}$, from which the optimal bandwidth is inferred. However, the bandwidth set does not have to be fixed in time; it can be dynamic of the same finite cardinality all the time. Namely, as time goes, all bandwidth $g$ values become useless except the one that takes all of the weight. This pattern can be traced and one can continuously inject/replace the obsolete candidate bandwidths with new larger ones such that the cardinality is kept fixed. Hence, our introduced method
is certainly still practical under such a continuously improving optimality.

Our convergence result is in line with the weighted majority approach in the mixture of experts [35], [40]. Briefly, we assign time-varying probabilities \(a_t^g(g)\) to each node on the tree for all bandwidth values in the set \(S\) such that \(\sum_{g \in S} a_t^g(g) = 1, a_t^g \geq 0, \forall t\). Initially, all bandwidths at a node \(\theta_k\) have equal priors, i.e., \(a_0^g(\theta_k) = \frac{1}{|S|}, \forall g \in S\). Then, we update these probabilities as in lines 10 and 12 in Algorithm 1 to converge to the best bandwidth in the set in the sense of likelihood maximization by learning the combination parameters \(a_t^g(\theta_k)\) in time. Namely, the final density estimation of the node \(\theta_k\) is a combination of density estimations evaluated with different bandwidth values as

\[
f_t^\theta_k(x) = \sum_{g \in S} a_t^g(\theta_k) f_t^g(x; g)
\]

where \(f_t^g(x; g)\) is the density estimation evaluated with bandwidth \(g\) as in (10) and \(f_t^\theta_k(x)\) is asymptotically tuned to the best \(g^*\) since \(a_t^g(g^*) \to 1\) as \(t \to \infty\). The reason is that we define \(a_t^g(\theta_k)\) through the likelihoods of the bandwidths based on the sum of their losses. Hence, the weight of the bandwidth value, which provides the best performance among the others, converges to 1, while all the others converge to 0 in time based on the learning parameter \(h\).

Note that this explanation for convergence to 1 is valid if the data statistics is stationary or asymptotically stationary by allowing nonstationarity in the transient state, which is necessary for the best performing \(g\) to consistently accumulate a performance that is increasingly better compared with others so that the best performing \(g\) can take all the weight. In case of nonstationarity (the distribution that the data instances are coming from might be changing in time, i.e., might be time-varying), the aforementioned convergence is not true but even in this case, we still can identify the best \(g\) by the largest weight.

Similarly, for a given partition \(P = \{\theta_1, \theta_2, \ldots, \theta_N\}\), we obtain bandwidth optimized piecewise online KDE as

\[
f_t^P(x) = \sum_{g \in S} a_t^g(\theta_k) f_t^g(x; g)
\]

We emphasize that since \(a_t^g(g^*) \to 1\) as \(t \to \infty\), our approach successfully identifies the best bandwidth in region \(\theta_k\) in a time-varying manner. The presented convergence result here indicates that we essentially favor the KDEs with smoother bandwidths in the beginning of the stream and gradually switch to the ones with steeper bandwidths as the data overwhelm, which is specifically to handle the overfitting issue directly from the bandwidth perspective.

**B. Efficient Combination of Hierarchically Organized Bandwidth Optimized Piecewise Online KDEs**

We point out that one can produce a bandwidth optimized piecewise online KDE for every given partition of the observation space. On the other hand, it is possible to obtain a doubly exponential large number, i.e., \((1.5)^{2D}\) with \(D\) being the depth, of many different partitions by all possible prunings of our partitioning tree. Hence, it is possible for one to approximate any desired partition with a tree-defined partition by using a relatively small \(D\).

It is critical to use the right partition in the case of anomaly detection with density estimation. Specifically, in the case of sequentially observed instances, the optimal partition is also not fixed but time-varying. Simpler partitions with a less number of regions are advantageous at the beginning of the stream, whereas one needs to use partitions with more complexities as more data become available. Hence, our algorithm produces a weighted combination of the piecewise online KDEs each of which corresponds to a specific partition. The combination weights are obtained based on the performance, i.e., the data likelihood, of each partition. Hence, the weights are time-varying and our algorithm relies mostly on the best-performing partitions due to the weighting where a better performing partition always receives a higher weight. In addition, we mathematically prove that our algorithm asymptotically performs at least as well as the best piecewise online KDE from the introduced large class of tree-defined partitions.

To this end, we introduce a computationally highly efficient combination of all of the bandwidth optimized piecewise online KDEs of prunings of our partitioning tree.\(^3\) The resulting final density estimator is given as

\[
f_t(x) = \sum_{k=1}^{n_p} c_t^k f_t^{P_k}(x) \tag{13}
\]

where \(n_p \approx (1.5)^{2D}\) and \(c_t^k\) is the weight of the partition \(P_k\) at time \(t\) as

\[
c_t^k = 2^{-\rho(P_k)} \exp \left(-h \sum_{u \in \theta_k, 1 \leq k \leq l} \left(\tilde{f}_t^{\theta_k∪}\right)\right) \tilde{C}
\]

with \(\tilde{C} = \sum_{l=1}^{(1.5)^{2D} - \rho(P_k)} \exp \left(-h \sum_{u \in \theta_k, 1 \leq k \leq l} \left(\tilde{f}_t^{\theta_k∪}\right)\right)\) the normalization constant, where \(\sum_{u \in \theta_k, 1 \leq k \leq l} \left(\tilde{f}_t^{\theta_k∪}\right)\) is the total loss of the partition \(P_k\) and \(\theta_k\) is the corresponding node of \(P_k\) at time \(u\). We define our loss for the node \(\theta_k\) as

\[
l_t^{\theta_k} = -\log (f_t^{\theta_k}(x)) \tag{14}
\]

which is the pointwise log-likelihood based on the estimator at that node. Also, \(\rho(P_k) = n_i + li - 1\) corresponds to the number of bits required to represent the partition \(P_k\), where \(n_i\) is the total number of nodes in partition \(P_k\) and \(l_i\) is the number of nodes with depth smaller than \(D\) in the partition.

On the contrary, combining \((1.5)^{2D}\) partitions or running \((1.5)^{2D}\) many bandwidth optimized piecewise online KDEs in parallel is computationally intractable. However, there exists an efficient solution [35], [36] to combine them with computational complexity only \(O(D)\). Instead of combining the partition specific estimations as in (13), one can obtain the same exact final estimation result of (13) by only combining \(D + 1\) many node specific estimations as

\[
f_t(x) = \sum_{k=0}^{D} c_t^k f_t^{g_k}(x) \tag{15}
\]

3For notational simplicity, we represent the normalized bandwidth optimized density estimation of the node \(\theta_k\) at time \(t\) as \(f_t^{\theta_k}(x)\) in the rest of this article by \(f_t^{\theta_k}(x) = \frac{\tilde{f}_t^{\theta_k}(x)}{\left(\tilde{f}_t^{\theta_k∪}\right)}\).
where $\tilde{f}^\theta_k(x)$ is the density estimation of the corresponding node at the $k$th layer of the tree and $c_{i}^\theta_k$ is the weight of the node $\theta_k$ at time $t$.

**Remark:** Recall that the density estimation of a partition is equal to the estimation of the corresponding node in that partition. Hence, although we have $(1.5)^2D$ many different partitions, we only have $D + 1$ unique estimations since the instance travels through only $D$ layers. Therefore, combining the partitions as in (13) exactly matches to combining the corresponding nodes at layers as in (15) with recollected weights in (16). This allows us to solve the problem in a more efficient way since it requires combining only $D + 1$ density estimations.

The weight $c_{i}^\theta_k$ of the node $\theta_k$ can be collected as [41]

$$
c_{i}^\theta_k = \sum_{P_k : j_f^T_k(x) = j_i^T_k(x)} c_{i}^{P_k}.
$$

(16)

The weight $c_{i}^\theta_k$ can also be recursively calculated by the evaluated loss or reward of the nodes as explicitly shown in [36]. We directly provide the recursion here while referring the reader to [36] for the details of the proof of this recursion. Then, one can obtain the weight of the node $\theta_k$ as

$$
c_{i}^\theta_k = \frac{2^{-\sum_{j=1}^{D_k} \rho(h)} \times \prod_{j=1}^{D_k} L_{\theta_j}(x^{t-1})}{L_{\theta_0}(x^{t-1})} \times \exp \left(-h \sum_{u : x_u \in \theta_k, 1 \leq u < t} l_u^\theta\right)
$$

(17)

where $\theta_0$ is the root node, $D_k$ is the layer of $\theta_k$ on the tree, $\theta_j$ is the other child of the parent of $\theta_k$ at the $j$th layer of the tree, and $L_{\theta_j}(x^{t-1})$ is an auxiliary variable that allows the recursive calculation [36] for any node $\theta_j$ using $L_{\theta_j}(x^{t-1})$

$$
\begin{align*}
&= \begin{cases} 
\exp \left(-h \sum_{u : x_u \in \theta_j, 1 \leq u < t} l_u^\theta\right), & (\theta_j \text{ is leaf}) \\
\frac{1}{2} L_{\theta_j}(x^{t-1}) L_{\theta_j}(x^{t-1}) + \frac{1}{2} \exp \left(-h \sum_{u : x_u \in \theta_j, 1 \leq u < t} l_u^\theta\right), & (\text{otherwise})
\end{cases} \\
\end{align*}
$$

(18)

with $\theta_l$ and $\theta_r$ being the left and right child nodes of the node $\theta_j$, respectively. We point out that this recursion is of computational complexity $O(D)$.

Based on this recursion (18) (see line 13 in Algorithm 1) and the description in (15) (see lines 7 and 8 in Algorithm 1), which is based on the bandwidth probability assignments (12) (see lines 7, 10, and 12 in Algorithm 1), we obtain our computationally efficient algorithm presented in Algorithm 1.

The following theorem provides the bound for the regret of our final density estimator (2), which combines the density estimations of the partitions.

**Theorem 1:** When our KDE-based hierarchical model is applied with Algorithm 1, the regret defined in (2) is upper bounded as

$$
R(T) = \sum_{t=1}^{T} (-\log(f_t(x_t))) - \min_{\theta \in \mathcal{P}} \left\{ \sum_{t=1}^{T} (-\log(f^\theta_t(x_t))) \right\}
$$

$$
\leq \frac{\rho(\mathcal{P}_i) \log(2)}{h} \leq \left(\frac{2D-1}{h}\right) \log(2) \forall i,
$$

(19)

where $\mathcal{P}_i$ is the $i$th partition in the competition class of estimators, $T$ is the length of sequence $x^T = [x_t]_{t=1}^{T}$, $h$ is the learning rate, and $\rho(\mathcal{P}_i) = |n_i + l_i - 1|$ measures the complexity of each partition by the number of bits required to represent that partition $\mathcal{P}_i$. Here, $\rho(\mathcal{P}_i)$ is defined based on the total number $n_i$ of nodes in the partition and the number $l_i$ of nodes with depth smaller than $D$ in the partition.

**Proof of Theorem 1:** The proof follows similar lines to the one in [35], with the difference that we use the negative log loss of the pointwise likelihood instead of the squared error loss. Following the definition in (18), $L_{\theta_0}(x^T)$ for the root node can be written as [36]:

$$
L_{\theta_0}(x^T) = \sum_{i=1}^{(1.5)^2D} 2^{-\rho(\mathcal{P}_i)} \exp \left(-h \sum_{t=1}^{T} l_t^\theta\right)
$$

where $\theta_i$ is the corresponding node of partition $\mathcal{P}_i$ at time $t$ and $l_t^\theta = -\log(f^\theta_t(x))$ as in (14). Then, we obtain

$$
L_{\theta_0}(x^T) \geq 2^{-\rho(\mathcal{P}_i)} \exp \left(-h \sum_{t=1}^{T} l_t^\theta\right)
$$

(20)

Also, when we apply (18) recursively to the root node, we observe that

$$
\frac{L_{\theta_0}(x^T)}{L_{\theta_0}(x^{t-1})} = \sum_{k=0}^{D} c_{i}^\theta_k \exp \left(h \log \tilde{f}^\theta_k(x_t)\right)
$$

where $c_{i}^\theta_k$ is defined as in (17). Since $\sum_{k=0}^{D} c_{i}^\theta_k = 1$ and $\exp(h \log(\cdot))$ is concave for the learning rate $0 \leq h \leq 1$, we obtain by Jensen’s inequality

$$
\frac{L_{\theta_0}(x^T)}{L_{\theta_0}(x^{t-1})} \leq \exp \left(h \log \left(\sum_{k=0}^{D} c_{i}^\theta_k \tilde{f}^\theta_k(x_t)\right)\right)
$$

where $f_t(x_t) = \sum_{k=0}^{D} c_{i}^\theta_k \tilde{f}^\theta_k(x_t)$ is the final density estimation (15). Since we have

$$
L_{\theta_0}(x^T) = \prod_{t=1}^{T} \frac{L_{\theta_0}(x^T)}{L_{\theta_0}(x^{t-1})}
$$

$$
\leq \prod_{t=1}^{T} \exp(h \log(f_t(x_t))) = \exp \left(h \sum_{t=1}^{T} \log(f_t(x_t))\right)
$$
where \( L_0(x^0) = 1 \), we write (20) as

\[
\exp \left( \sum_{t=1}^{T} \log (f_t(x_t)) \right) \geq L_0(x^T)
\]

\[
\geq 2^{-\rho(P_t)} \exp \left( -h \sum_{t=1}^{T} \beta_t \right)
\]

\[
\frac{1}{h} \sum_{t=1}^{T} (- \log (f_t(x_t))) \leq \frac{\rho(P_t) \log(2)}{h}
\]

Finally, we obtain

\[
\frac{1}{h} \sum_{t=1}^{T} (- \log (f_t(x_t))) - \sum_{t=1}^{T} (- \log (f_t(x_t))) \leq \frac{\rho(P_t) \log(2)}{h}
\]

Since \( \rho(P_t) \leq 2^D - 1 \), we obtain the following regret bound

\[
R_T = \frac{\rho(P_t) \log(2)}{h}
\]

Hence, the normalized regret \( R_T/T \) is convergent to 0 at a \( O(1/T) \) rate. This shows that our density estimation algorithm is regret optimal, i.e., it asymptotically performs (in terms of the accumulated log likelihood) as well as the piecewise online KDE that corresponds to the best tree-defined partition.

V. ANOMALY DETECTOR

We compare the estimated probability density \( f_t(x_t) \) at each time \( t \) with a threshold to decide whether \( x_t \) is anomalous. A data instance \( x_t \) having a sufficiently low probability density is considered anomalous with respect to the rule

\[
\hat{d}_t = \begin{cases} +1, & f_t(x_t) < \varsigma_t \text{ (anomalous)} \\ -1, & f_t(x_t) \geq \varsigma_t \text{ (normal)} \end{cases}
\]

Based on this rule, we find the optimal threshold that achieves the best anomaly detection performance by updating its value in time. For this, we use the logistic loss function of \( \varsigma_t \) that is to be minimized [42]:

\[
r_t = \log(1 + \exp(- (\varsigma_t - f_t(x_t))d_t)).
\]

Note that the loss is small if \( d_t = \hat{d}_t \), and it increases if the difference \( (\varsigma_t - f_t(x_t)) \) is increasing when \( d_t = \hat{d}_t \). Hence, the model learns the optimal threshold that differentiates the normal and anomalous data in time with loss minimization.

Then, the threshold is updated by using the OGD method [11], which learns the optimal parameter minimizing the desired loss, as in Algorithm 2. We calculate the gradient of the loss function with respect to \( \varsigma_t \) as \( \nabla_{\varsigma_t} r_t = -d_t (1 + \exp((\varsigma_t - f_t(x_t))d_t))^{-1} \) and the update follows as \( \varsigma_{t+1} = \varsigma_t - \eta_t \nabla_{\varsigma_t} r_t \), where \( \eta_t \) is the learning rate, which depends on the number of data instances and the convexity of the defined loss function [11], and \( P_r(x) = \arg \min_{y \in \mathcal{T}} ||x - y||^2 \) is the projection onto convex threshold set \( \mathcal{T} \). Since the loss function \( r_t \) is convex, our algorithm converges to the optimal threshold in the set by minimizing the cumulative loss.

VI. EXPERIMENTS

In this section, we extensively evaluate the performance of our anomaly detector on an artificial data set of a mixture of two Gaussian components, and several other real data sets, such as Occupancy [43], Breast Cancer Wisconsin [44], Pen digits [44], Susy [45], Thyroid [44], Pima [44], Mammography, and Liver [44]. Comparisons with one-class SVM [24], incrementally trained SVM (we represent as “I-SVM” throughout this section) [46], K-nearest neighbors (K-NNs) [25], K-D tree nearest neighbor search (K-D Tree) [47], K-localized p-value estimation (K-LPE) [9], online oversampling PCA (online osPCA) [48], and Fourier OGD (FOGD) [34] are presented, which are commonly used for anomaly detection. We also include comparisons with anomaly detection method based on KDE with local bandwidths where the bandwidth depends on the distance to the nearest neighbor of the observed point as in [49] and represent this technique as “K-KDE.” Finally, we include the performance of our method without partitioning, which is represented as “kernel density estimation with adaptive bandwidth (KDE-AB).” We use the receiver operating characteristics (ROC) and the area under curve (AUC) to compare the anomaly detection performances of the algorithms. The ROC curves are obtained by plotting the rates of true detection versus false alarm while varying the corresponding threshold in each algorithm. For this, we use the scores, i.e., values that are thresholded to detect anomalies, assigned to the observed instances by the algorithms except for K-NN and K-D Tree. For them, we consider the score as the proportion of the neighbors with positive class label. The parameters \( (\nu \text{ and } \gamma) \) for the SVM and I-SVM and \( K \) for the K-NN, K-D Tree, and K-LPE of the compared algorithms are optimized with tenfold cross validation, where we reserve the 75% of each data set for training and the remaining for the test. Since the other methods, including our algorithm AD-HKDE, does not need cross validation as they are online, we set their parameters, i.e., \( D, h, \) and \( G \) for AD-HKDE, based on a small preliminary experiment. In Algorithm 1, the tree depth and the learning rate are fixed as \( D = 3 \) and \( h = 0.01 \), respectively.

Since, in general, the higher layer nodes on our partitioning tree observe an exponentially smaller number of data instances compared to the lower layer ones, we form the bandwidth set \( G = \{g_1, \ldots, g_k\} \) with relatively small \( g \) values for the root node and larger \( g \) values for the leaf nodes. We use \( |G| = k = 4 \) and form the sets such that \( G = \mathcal{E} \subset \{2^0, 2^1, \ldots, 2^{k-1}\} \), where we increase \( \epsilon = \mathbb{L}(l) \) (\( l \) denotes

\[\text{Algorithm 2 Anomaly Detector}\]

1. Specify \( \eta_t \) and initialize the learning rate \( \eta_t = \frac{1}{\eta} \)
2. Specify the threshold set \( T \) and initialize \( \varsigma_t \)
3. for \( t = 1, 2, \ldots \) do
4. Calculate the learning rate \( \eta_t = \frac{1}{\eta} \)
5. Compare \( f_t(x_t) \) with \( \varsigma_t \) and declare the decision \( \hat{d}_t \)
6. Observe the actual \( d_t \) and calculate \( \nabla_{\varsigma_t} r_t \)
7. Update the threshold \( \varsigma_{t+1} = \varsigma_t - \eta_t \nabla_{\varsigma_t} r_t \)
8. end for
the layer and $L = \{0.01, 0.5, 1.5, 2\}$ from the root to the leaves. Hence, the union of all such sets from the layers of the tree yields a relatively large actual set (whose size is proportional to the tree depth) of bandwidth parameter values in which our algorithm optimizes for $g$ with low computational and space complexity. Note that, in general, optimal kernel bandwidth $\delta$ is proportional to $N^{-(1/5)}$ ($N$ is the number of data points) [12] ($g = (1/\delta)$ is proportional to $N^{(1/5)}$). In the online framework, the number of data points changes from $1$ to $T$ ($T$ is the total number of data instances). Therefore, in our experiments, we define a bandwidth set such that its maximum is proportional to $T^{(1/5)}$ based on our largest data set. Finally, the dimensionality is reduced to $d = 3$ for all data sets (except the artificial one which is already 2-D) by using PCA for only partitioning purposes; otherwise, the full dimension is used in the remaining parts such as estimating the density. In Algorithm 2, the learning rate parameter is set as $\eta = 0.001 \in (0, 1)$ and the threshold set is $T = [10^{-5}, 2 \times 10^{-5}, \ldots, 1]$.

In this section, “AD-HKDE” represents our technique.

A. Artificial Data Set

In the first part of our experiments, we generate a 2-D data set of length $T = 2000$, whose probability density is a mixture of two Gaussian components $f \sim (1/2)N \begin{bmatrix} -1 & 1 \\ -1 & 0 \end{bmatrix} + (1/2)N \begin{bmatrix} 1 \\ 1 \end{bmatrix} \cdot \begin{bmatrix} 4 & 0 \\ 0 & 8 \end{bmatrix}$ as the normal data. Our algorithm does not need anomalous data to work; however, we replace the normal data at 400 randomly picked instances with uniformly distributed anomalous observations to test the anomaly detection performance.

As can be seen in Fig. 2(a), our algorithm AD-HKDE outperforms all the others in terms of the AUC score while being largely superior in either relatively smaller or higher false alarm rate regions. We point out to understand the reason that the data consist of two Gaussian components, where the second component has four or two times larger variance in its dimensions compared with the other component. This specifically detrimentally affects the methods K-NN, K-D Tree, and K-LPE. The reason is that, in regions of the higher variance component, one has to base the estimations on wider neighborhoods (compared with the other component) in order to observe the same $k$ number of instances. This certainly degrades their estimation since the underlying density is probably far from being uniform in such wide neighborhoods. K-KDE suffers from the placements of the data points in the observation space due to its pointwise bandwidth. Since its kernel bandwidth is directly proportional to the distance between the observed instance and its nearest neighbor, it is not able to learn the optimal bandwidth and estimate the density correctly. Online osPCA uses only the first principal direction, and hence, it loses the information carried by the second component, which decreases its performance significantly. FOGD, however, suffers from the fixed kernel bandwidth, which does not allow the algorithm to learn the local variations. For the same reason, our algorithm AD-HKDE outperforms SVM and I-SVM since they fit symmetric RBF kernels with equal bandwidths at each instance for each component regardless of the component covariances and the time-varying data size. In accordance, our algorithm exploits the local structure of the density and learns the optimal bandwidths for each component even in a time-varying manner. Therefore, our algorithm AD-HKDE also outperforms its nonpartitioned version, i.e., KDE-AB, which learns the bandwidth based on the whole observation space.

We emphasize that the algorithms SVM, incrementally trained SVM, K-NN, K-D Tree, and K-LPE with separate training and test phases, are specifically chosen in our experiments due to their competing high performance and comparable model-free working environment, but their computational complexity is significantly larger than the one of our online detector. For example, SVM has complexity $O(T_d)$ in training, and $O(T_d T_{test})$ in test phase (both are in the worst case). Incrementally trained SVM aims to reduce the computational complexity by using a time-varying manner. Therefore, our algorithm AD-HKDE also outperforms its nonpartitioned version, i.e., KDE-AB, which learns the bandwidth based on the whole observation space.

![ROC Curves](image1)

Fig. 2. ROC performance of the compared algorithms on (a) mixture of two Gaussian components and (b) Occupancy data set.
complexity of SVM. This method proposes to divide the training data into $N$ subsets, which has complexity $O(T_d^2/N^2)$ in training and $O((T_d T_{test})/N)$ in test phase. The complexity of K-LPE is approximately $O(T_d^2)$ in training and $O(T_d T_{test})$ in test phase. The complexity of K-NN is $O(T_d T_{test})$, which is prohibitively costly when the number of data points is high. K-D Tree decreases the computational complexity of K-NN by introducing a tree structure. This method has complexity $O(T_d \log(T_d))$ in training and $O(T_{test} \log(T_d))$ in test phase, which is still high for a large number of data points. On the other hand, the complexity of our detector is only $O(DT_k)$ ($T$ is the data length, $D$ is the depth of the partitioning tree, and $k$ is the number of $g$ values in the bandwidth set). Moreover, our algorithm is truly online without separate training or test phases, where increments are straightforward with new data; however, increments in these other competing ones are computationally demanding.

B. Real Data Sets

In this section, we present the ROC performances of the compared algorithms on several real data sets in Table I that are widely used in the anomaly detection literature. We first perform an experiment on the Occupancy data set [43]. This data set consists of 10,808 data points where the labels correspond to occupied (normal) and unoccupied (anomalous) room states.

As shown in Fig. 2(b), our technique AD-HKDE achieves the highest AUC score, which indicates that the intrinsic covariance structure within this data set heavily requires local as well as temporal bandwidth adaptation; therefore, the superiority of our algorithm follows. In all of our experiments including the Occupancy data set and the others, our algorithm AD-HKDE is generally highly superior in relatively smaller false alarm regions (except a few cases). Small false alarm rate requires the estimation of low density, which is in turn largely susceptible to insufficient data in the corresponding low-density regions. This necessitates to adjust the bandwidths or the neighborhood sizes intelligently in those regions of the support. Our algorithm elegantly addresses this need, which introduces its superiority.

In addition, we plot the temporal variations of the bandwidth probabilities/weights of our algorithm in Fig. 3(a) to show the bandwidth adaptation power of our method. As an example, we specify the bandwidth set as $G = \{0.01, 0.02, 0.04, 0.08\}$ for the root node $\theta_r$ and apply our algorithm on the Occupancy data set. Since there are four possible bandwidths, initial probabilities are $\alpha_{\theta_k}(g) = 1/4$, $\forall g \in G$. Then, as shown in Fig. 3(a), our algorithm asymptotically chooses $g = 0.08$ (the smallest bandwidth) since its probability converges to one, whereas the others converge to zero. This shows the efficacy of our algorithm since KDEs generally achieve higher performance with small bandwidths when the number of instances increases.

In Fig. 3(b), we provide the ROC of our algorithm and its static version with fixed bandwidths, i.e., with singleton $G$’s at the nodes. The number of data instances is low at the beginning

TABLE I

PROPERTIES OF THE REAL DATA SETS USED IN THE EXPERIMENTS

<table>
<thead>
<tr>
<th>Datasets</th>
<th>Properties</th>
<th>Number of features</th>
<th>Total number of instances</th>
<th>Number of normal instances</th>
<th>Number of anomalous instances</th>
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<td>8107</td>
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<td>569</td>
<td>357</td>
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<td>6714</td>
<td>156</td>
<td></td>
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<td>$54 \times 10^4$</td>
<td>$46 \times 10^4$</td>
<td></td>
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</table>

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As we define the bandwidth parameter as $g = (1/\sqrt{\delta^2})$, $g = 0.08$ corresponds to the actual bandwidth $\delta = 2.5$. 

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TABLE II

AUC Scores of the Algorithms for the Occupancy, Breast Wisconsin, Pen Digits, Susy, Thyroid, Pima, Bupa Liver, and Mammography Data Sets

<table>
<thead>
<tr>
<th>Datasets</th>
<th>SVM</th>
<th>I-SVM</th>
<th>K-NN</th>
<th>K-D Tree</th>
<th>Online osPCA</th>
<th>K-LPE</th>
<th>FOGD</th>
<th>K-KDE</th>
<th>KDE-AB</th>
<th>AD-HKDE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Occupancy</td>
<td>0.9828</td>
<td>0.9750</td>
<td>0.9854</td>
<td>0.9854</td>
<td>0.9292</td>
<td>0.9705</td>
<td>0.9490</td>
<td>0.9368</td>
<td>0.9531</td>
<td>0.9907</td>
</tr>
<tr>
<td>Breast Wisconsin</td>
<td>0.9483</td>
<td>0.9480</td>
<td>0.9778</td>
<td>0.9778</td>
<td>0.7248</td>
<td>0.9405</td>
<td>0.9555</td>
<td>0.8363</td>
<td>0.9083</td>
<td>0.9672</td>
</tr>
<tr>
<td>Pen digits</td>
<td>0.9362</td>
<td>0.9334</td>
<td>0.9858</td>
<td>0.9858</td>
<td>0.8951</td>
<td>0.9902</td>
<td>0.9503</td>
<td>0.9145</td>
<td>0.9769</td>
<td>0.9913</td>
</tr>
<tr>
<td>Susy</td>
<td>0.6821</td>
<td>0.6343</td>
<td>0.8569</td>
<td>0.8569</td>
<td>0.5747</td>
<td>0.7651</td>
<td>0.8474</td>
<td>0.6997</td>
<td>0.6928</td>
<td>0.9078</td>
</tr>
<tr>
<td>Thyroid</td>
<td>0.8809</td>
<td>0.8747</td>
<td>0.9208</td>
<td>0.9208</td>
<td>0.8772</td>
<td>0.9286</td>
<td>0.8937</td>
<td>0.7431</td>
<td>0.8734</td>
<td>0.9401</td>
</tr>
<tr>
<td>Pima</td>
<td>0.6694</td>
<td>0.6694</td>
<td>0.7690</td>
<td>0.7690</td>
<td>0.6272</td>
<td>0.7032</td>
<td>0.6309</td>
<td>0.6473</td>
<td>0.6652</td>
<td>0.7932</td>
</tr>
<tr>
<td>Liver</td>
<td>0.5950</td>
<td>0.5950</td>
<td>0.7062</td>
<td>0.7062</td>
<td>0.5631</td>
<td>0.6923</td>
<td>0.6733</td>
<td>0.7013</td>
<td>0.5756</td>
<td>0.7436</td>
</tr>
<tr>
<td>Mammography</td>
<td>0.8337</td>
<td>0.8295</td>
<td>0.8676</td>
<td>0.8676</td>
<td>0.7667</td>
<td>0.8709</td>
<td>0.8279</td>
<td>0.7058</td>
<td>0.8568</td>
<td>0.9115</td>
</tr>
</tbody>
</table>

and then increases since the algorithms operate in the online framework. Therefore, KDEs with smaller bandwidths (larger \( g \)) achieve higher bandwidth probabilities in Fig. 3(a), toward the end of the process. However, KDEs with fixed smaller bandwidths get detrimentally affected by the scarce data in the beginning in terms of the detection performance, which results in an overall lower ROC performance as in Fig. 3(b). For this contrast between Fig. 3(a) and (b), we point out that an improvement in the density estimation does not directly translate to an improvement in the ROC performance, as the ranking of point density estimates matters more in our detectors compared to raw density estimates. Also, Fig. 3(a) shows the dynamic bandwidth probabilities with respect to the changing data size, whereas Fig. 3(b) shows the overall AUC performance. Namely, a relatively larger bandwidth can still yield a better AUC in overall, although a smaller one is expected to perform better temporally locally toward the end of processing. Finally, our algorithm outperforms the static versions and has the highest AUC since it dynamically adapts its bandwidths in accordance with the available data size.

We also apply the algorithms on the Breast Cancer Wisconsin data set [44]. As shown in Table II, K-NN and K-D Tree are the best-performing ones in this case (K-D Tree is K-NN with lower computational complexity), and our method AD-HKDE is the second best (the other methods are strongly outperformed). The reason is that the data size is not large enough to perfectly learn the bandwidths in all regions across time. Namely, in our method, we require to learn more parameters (compared with K-NN and K-D Tree) and hence demand more instances. Even when the data size is not sufficient as in this case, the difference is so small, as shown in Table II. Also, note that our method works sequentially in the truly online setting with no separate training or test phases or even cross validation, while K-NN and K-D Tree need training data making them impractical in our processing setting. In our targeted data streaming applications, the data size is typically not an issue.

Moreover, we perform an experiment on pen digits data set [44] that is composed of handwritten digits from 0 to 9. Here, 0 digits are considered as anomalous, and digits from 1 to 9 with an equal number are considered as normal. This data set consists of ten digits, which can also be considered as the mixture of ten distributions. Since our technique AD-HKDE is able to localize these components piecewise in its density estimation with adaptive bandwidths, we obtain a decent performance on this data set with better AUC than all the others, as shown in Table II.

We also evaluate the AUC scores of the algorithms on the Thyroid [44], Pima [44], Mammography, and Liver [44] data sets, whose properties are given in Table I. As shown in Table II, our algorithm AD-HKDE outperforms the other methods in all cases with its high modeling capability. In particular, for the Thyroid and Mammography data sets, our method provides significantly higher AUC scores than the compared algorithms. Since the number of instances is greater in these data sets, our algorithm AD-HKDE learns the data structure more, which results in higher detection performance.

Finally, we perform an experiment on the Susy data set [45] with 10^6 instances. As shown in Table II, our algorithm AD-HKDE has the highest AUC score among all methods. In this case, we clearly observe the advantage of partitioning since our method AD-HKDE also significantly outperforms the other kernel-based methods where FOGD uses fixed bandwidth and K-KDE uses pointwise variable bandwidth based on the nearest neighbor distance. The number of data points being large allows our algorithm to perfectly learn the local structure of the data set and optimal partitioning with the corresponding bandwidths.

VII. Conclusion

We introduced an online unsupervised anomaly detection algorithm for sequential data. The probability density of the observed instance is first estimated based on the past observations. Our density estimation can be used even if the underlying distribution is of any complex form since we do not have any parametric shape assumption. After the density for a given data instance is estimated, an anomaly is declared if the estimated density is sufficiently low following a comparison to a data-adaptive threshold. Our density estimation is a mixture of experts approach based on an ensemble of hierarchically organized KDEs. The proposed algorithm guarantees to sequentially achieve the best likelihood performance within the introduced large ensemble represented by a binary partitioning tree. This mathematical performance guarantee implies that
our method essentially exploits fitting kernels of bandwidths that are optimally varying both in space and time in the KDE technique. Namely, our method learns the optimal partitioning of the observation space with KDEs trained separately in partition regions. Such partition region KDEs are even in partition regions. Such partition region KDEs are even.

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for processing at extremely fast rates with decent anomaly processing in our framework is truly online with computational complexity \(O(TDk)\) that is only linear in the tree depth \(D\), the length \(T\) of the data, and the number \(k\) of bandwidths in the bandwidth set. Hence, our technique is appropriate for processing at extremely fast rates with decent anomaly detection performance.

REFERENCES


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